

NCKU Programming Contest Training Course

Course 9

2015/03/25

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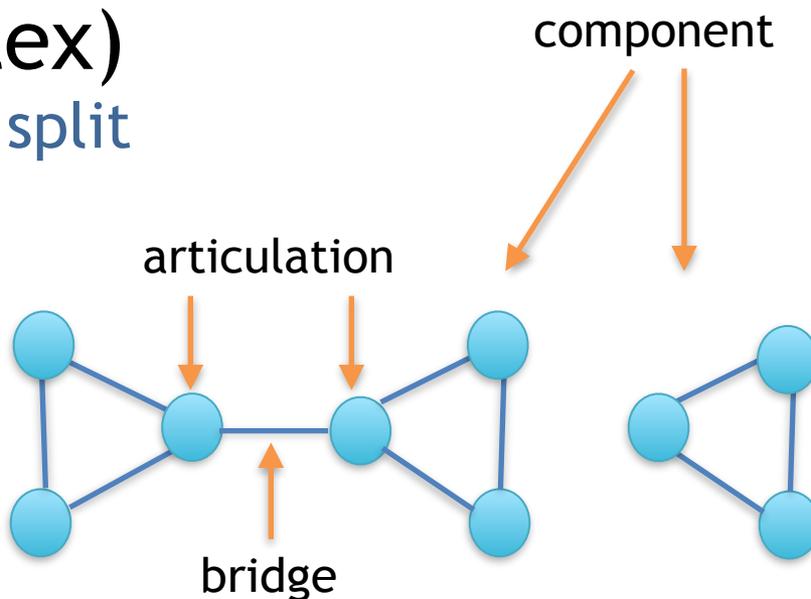
Outline

- Articulation/Bridge
- Strongly Connected Component(SCC)



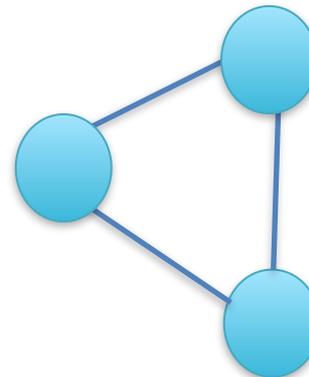
Graph

- **connected graph/component**
connected graph/component iff all pairwise vertices exist **at least one path** & **no more vertices can be added**
- **articulation(cut-vertex)**
remove articulation vertex split one component to two
- **bridge(cut-edge)**
same as articulation



Articulation / Bridge

- Find Articulation in Graph
 - Graph become **non-connected** if remove a Articulation.
V times DFS = $O(V*(V+E))$ → **too slow!**
 - Vertex is not Articulation if can find **alternative path**
→ find cycle!
 - Use DFS → $O(V+E)$

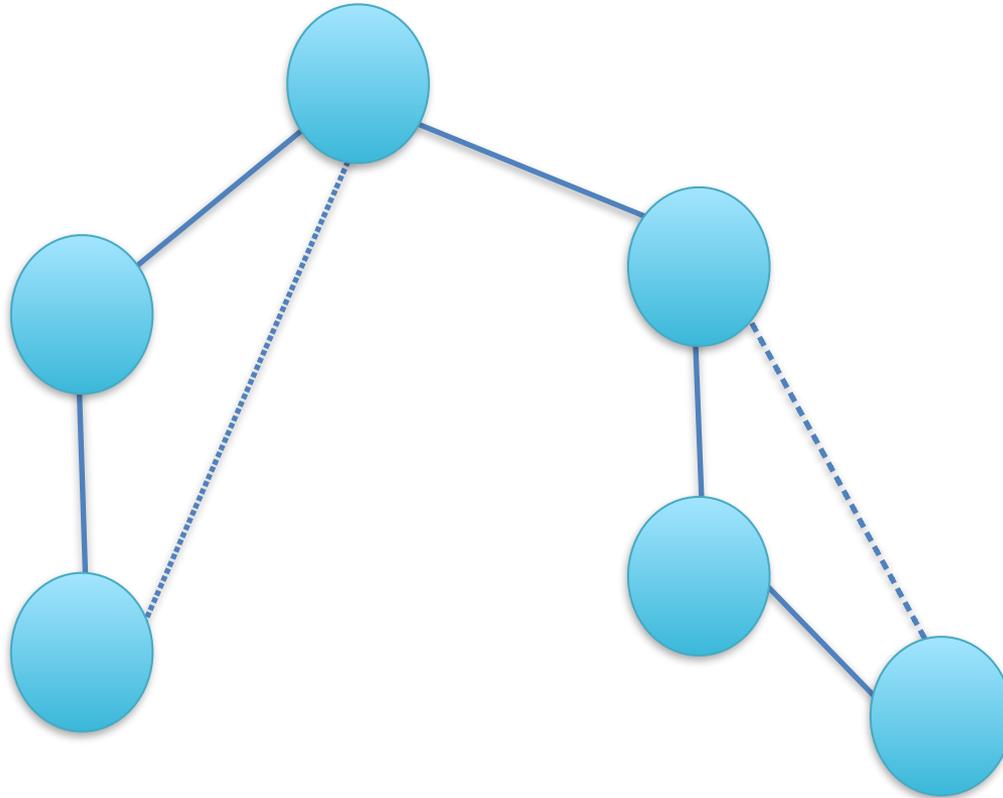


Articulation / Bridge

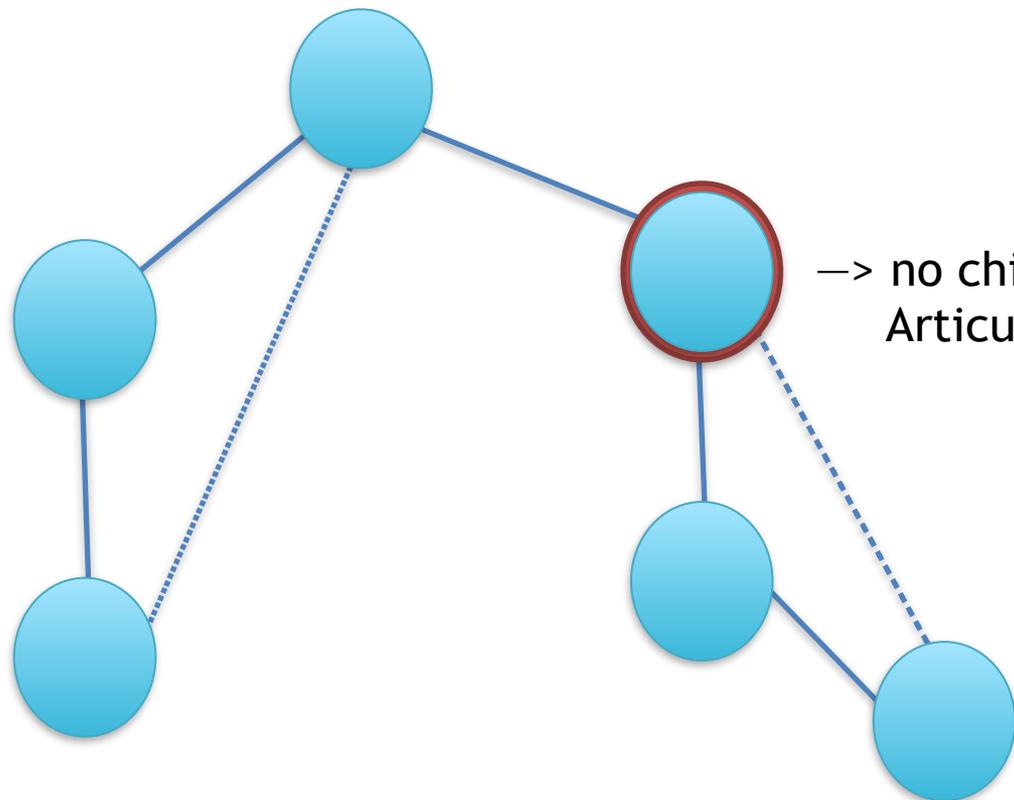
- Concept
 - if vertex u 's children can't back to u 's ancestors
→ u is Articulation
 - if vertex u is root and has at least 2 child
→ u is Articulation
- Bridge?
 - two Articulation u, v have an edge → (u, v) is Bridge!



Articulation/Bridge



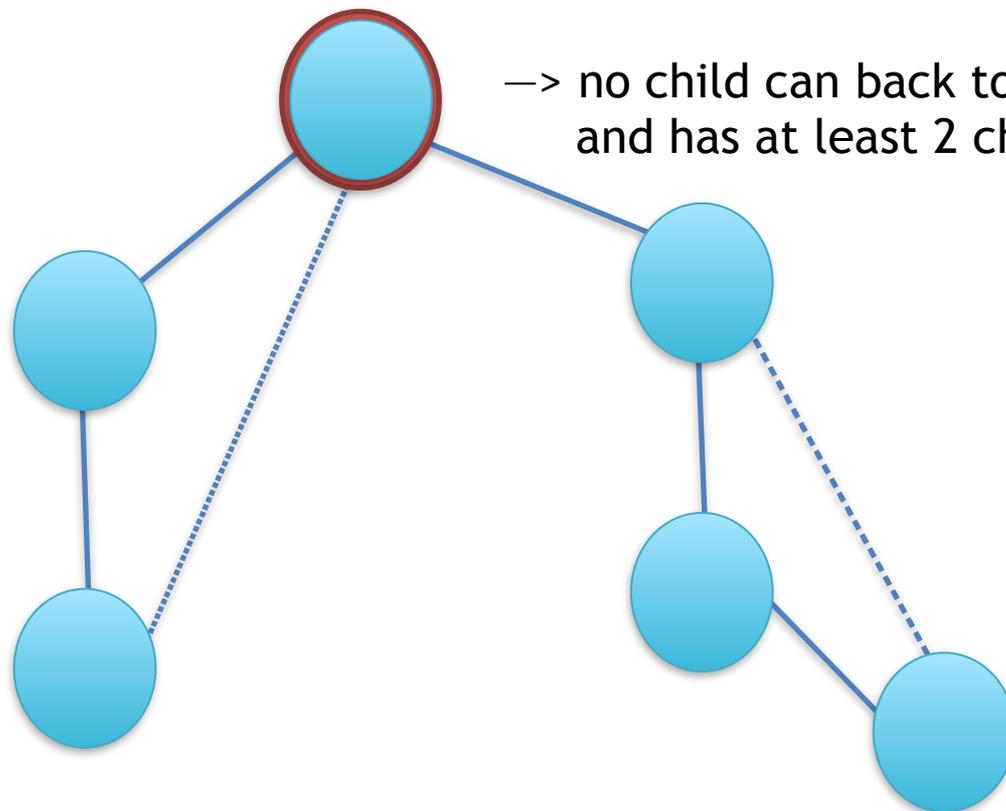
Articulation / Bridge



→ no child can back to ancestor.
Articulation!



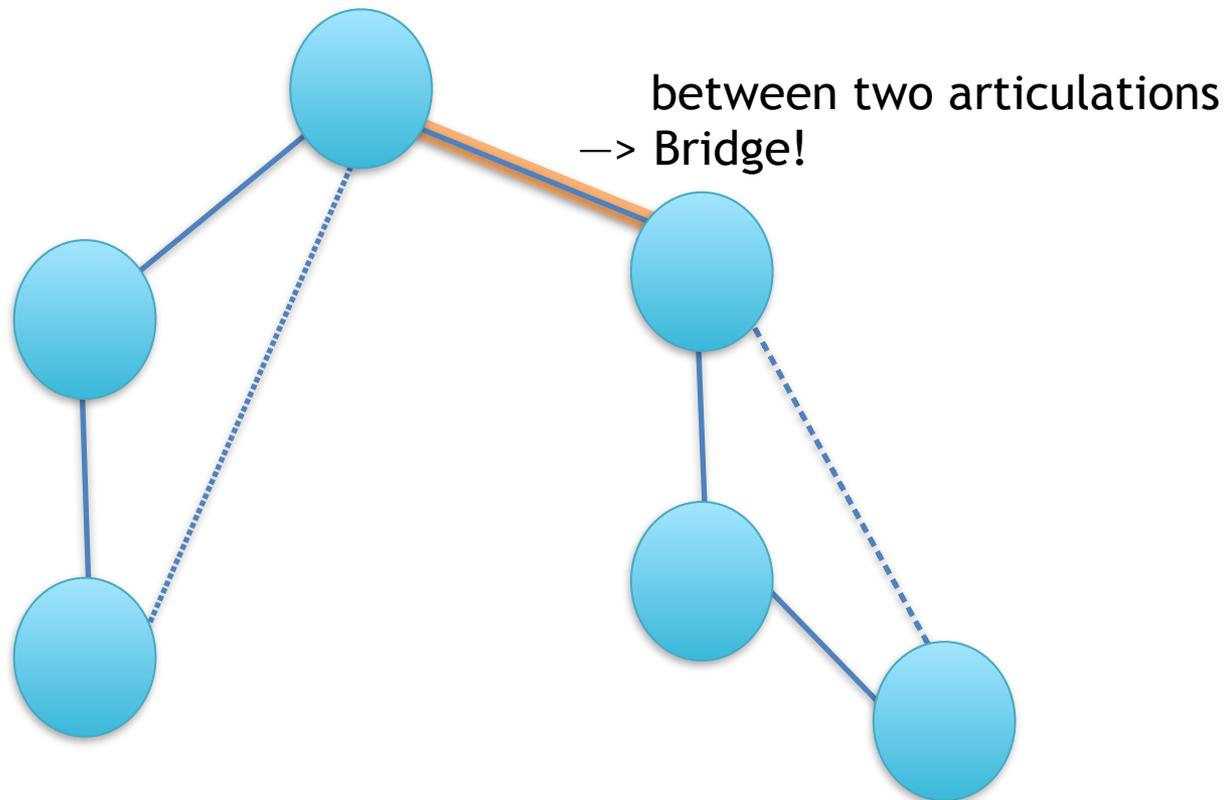
Articulation/Bridge



→ no child can back to ancestor and has at least 2 child. Articulation!

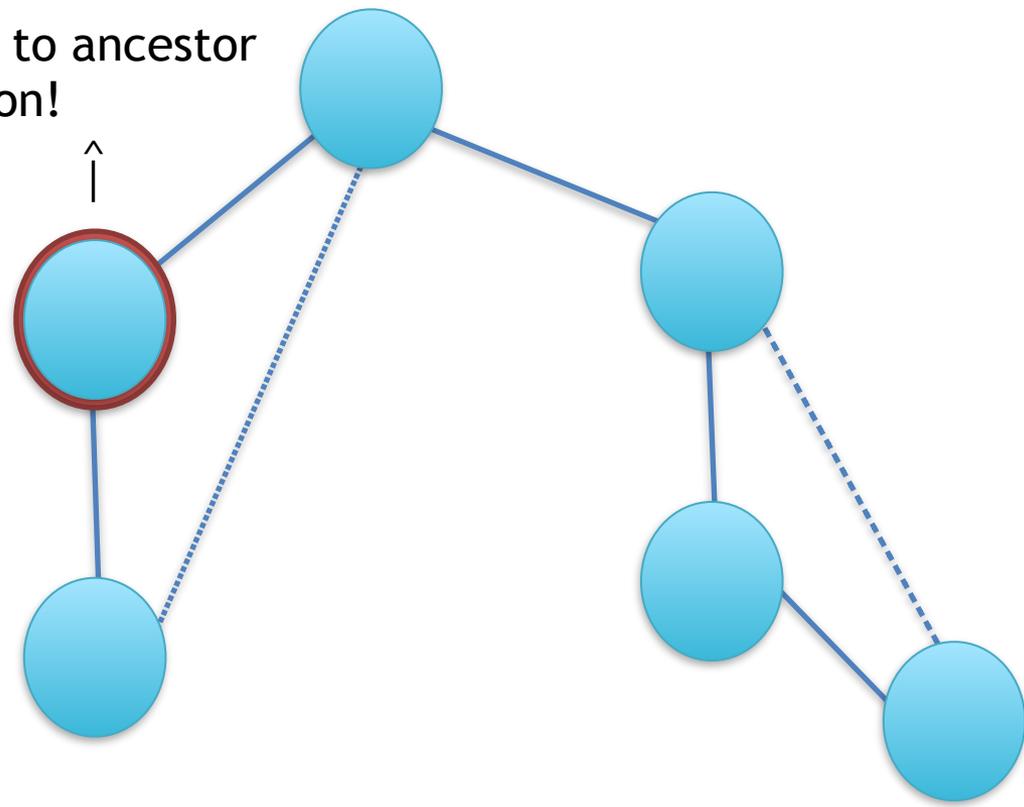


Articulation/Bridge



Articulation/Bridge

child can back to ancestor
NOT Articulation!



Articulation / Bridge

- $dfn[u]$ = DFS traversal order
 - first visit time each vertex u in DFS

- $low[u]$ = $\min(dfn[u], \text{lowest } low[v])$
 - if edge (u,v) exist and v is not u 's parent



Articulation / Bridge

- Articulation
 - if vertex u 's children can't back to u 's ancestors
→ $dfn[u] \leq low[v]$, v is u 's child
 - if vertex u is root and has at least 2 child
→ count $child \geq 2$
- Bridge?
 - two Articulation u, v → $dfn[u] < low[v]$, v is u 's child



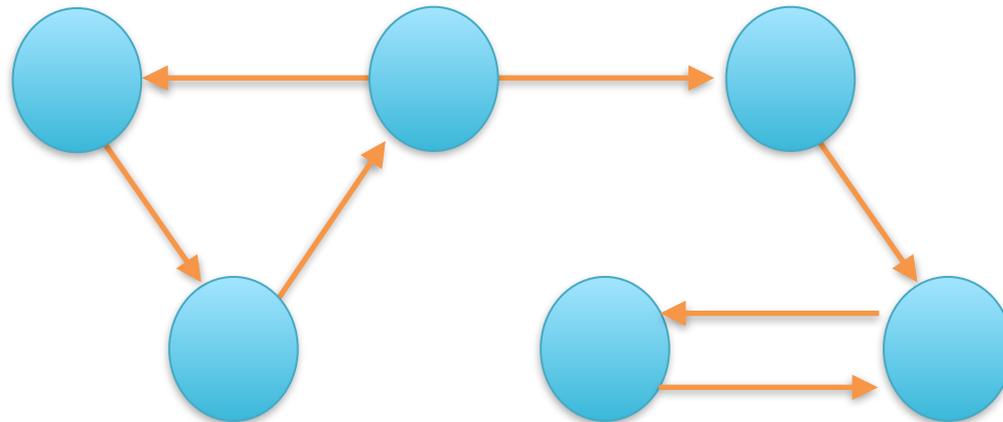
Practice

UVA - 315

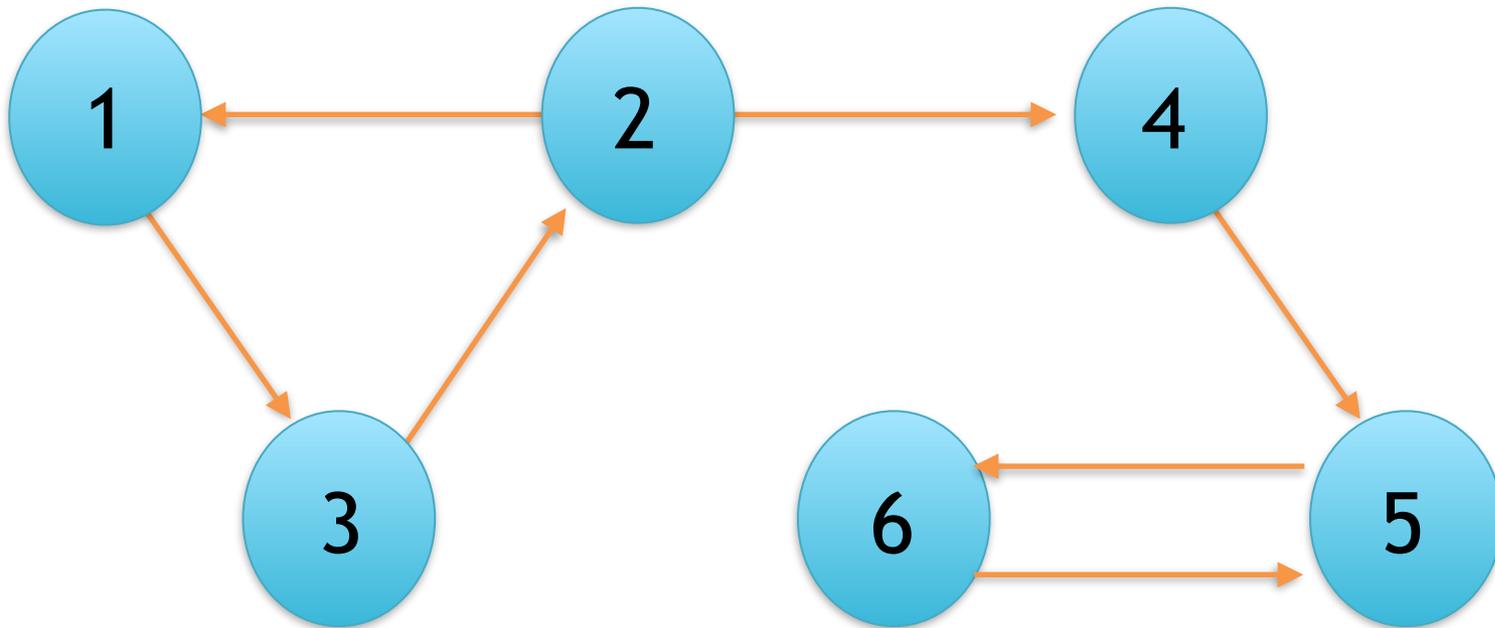


SCC

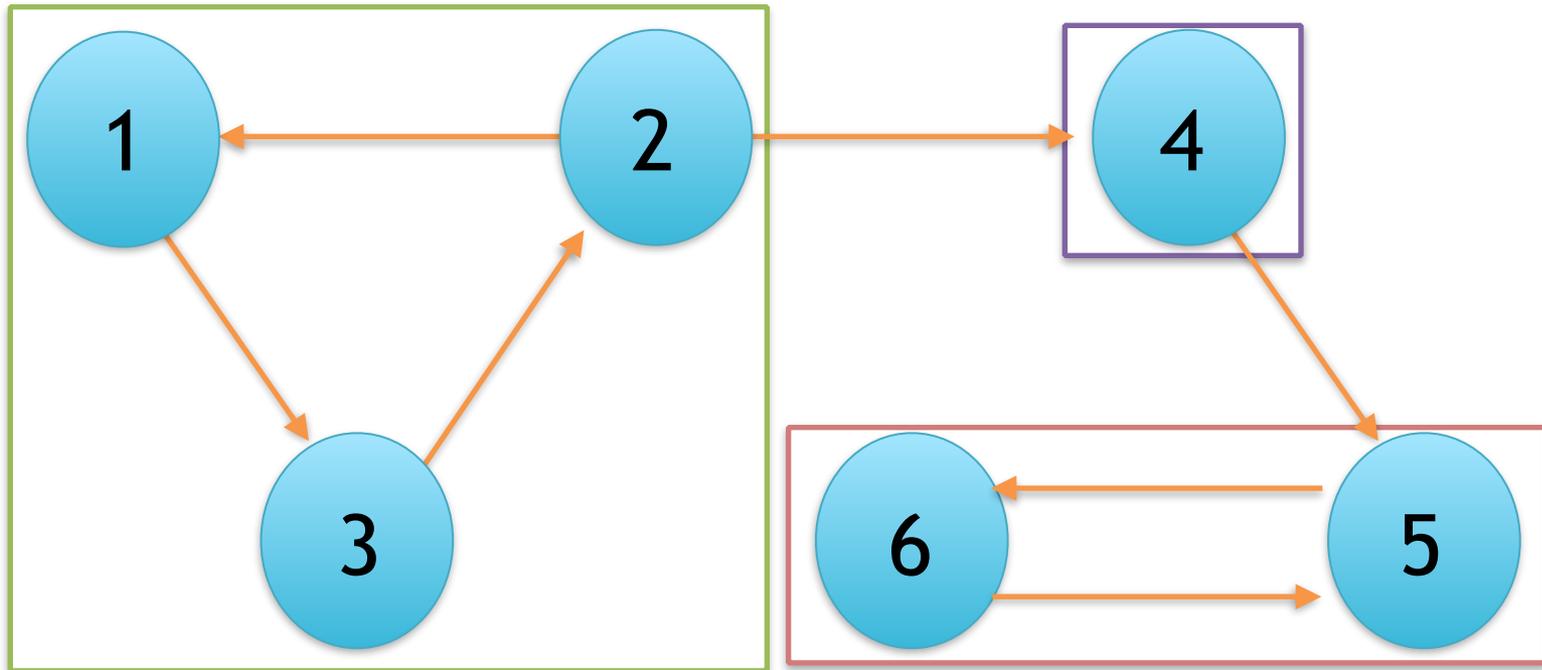
- connected component in **directed** graph
 - same definition in undirected graph



SCC



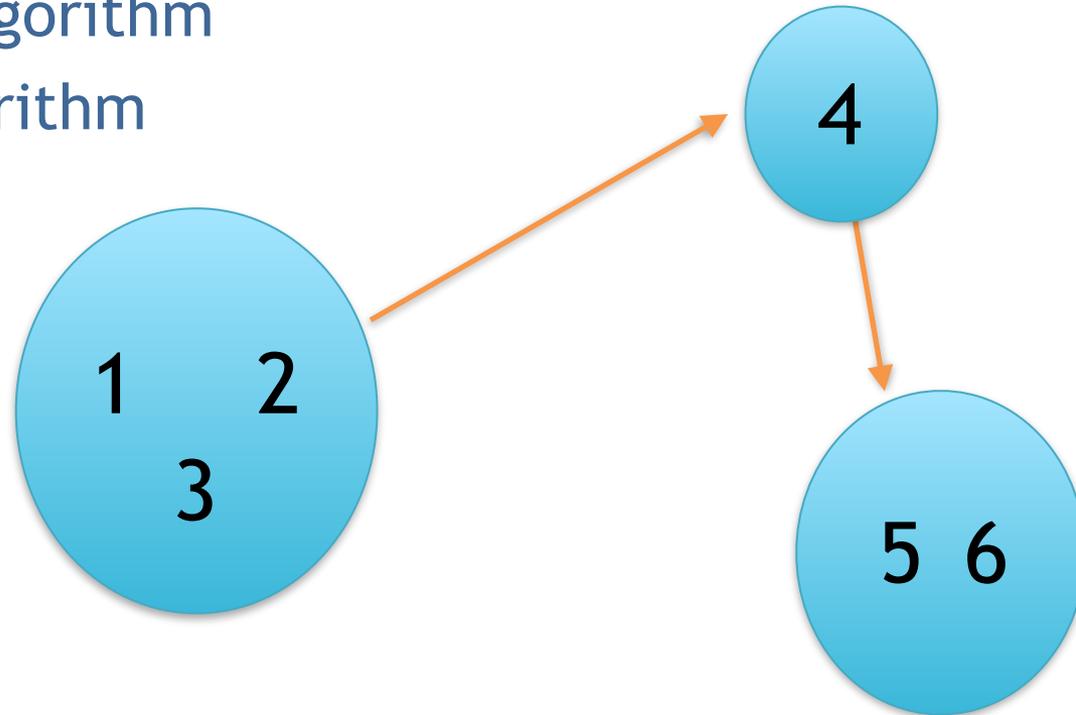
SCC



SCC

find all SCCs, **contract all cycles** → DAG (directed acyclic graph)

- Kosaraju's Algorithm
- Tarjan's Algorithm



SCC

- Kosaraju's algorithm

STRONGLY-CONNECTED-COMPONENTS(G)

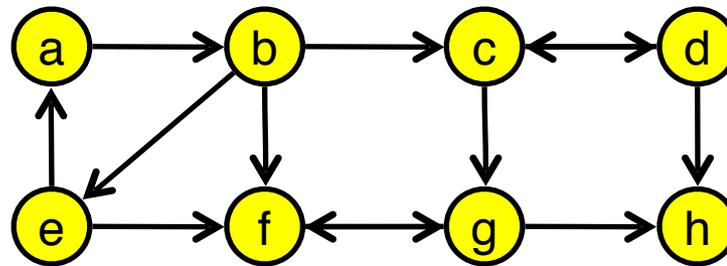
1. Call DFS(G) to compute finishing time for each vertex.
2. Compute transpose of G i.e., G^T .
3. Call DFS(G^T) but this time consider the vertices in order of decreasing finish time.
4. Out the vertices of each tree in DFS-forest.

twice DFS \rightarrow total complexity: $O(V+E)$



SCC

- Algorithm

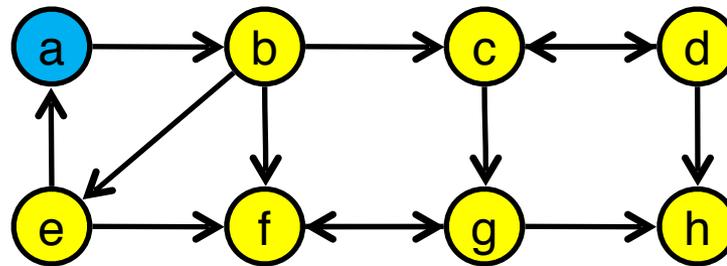


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SCC

- Algorithm

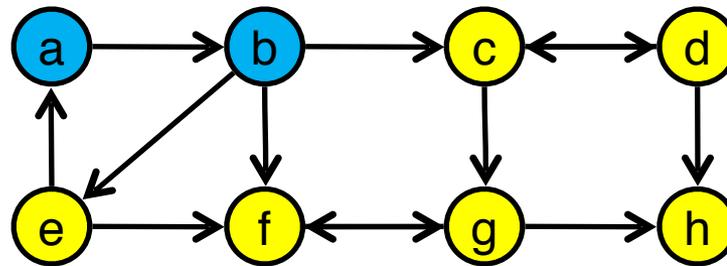


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SCC

- Algorithm

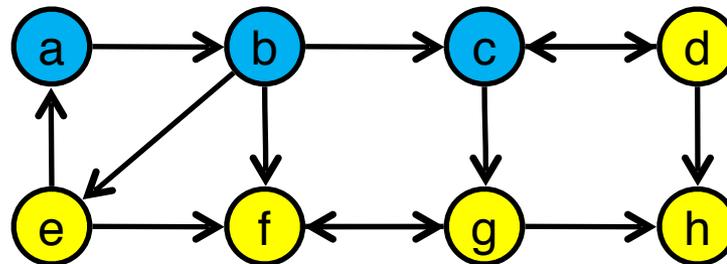


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SCC

- Algorithm

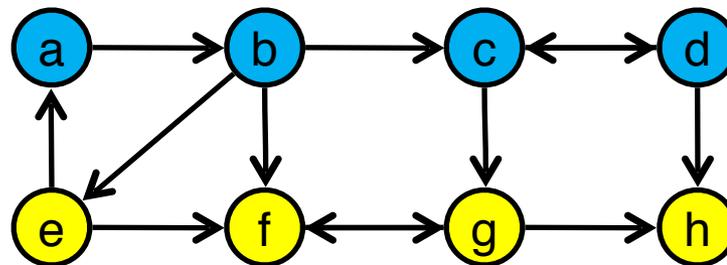


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SCC

- Algorithm

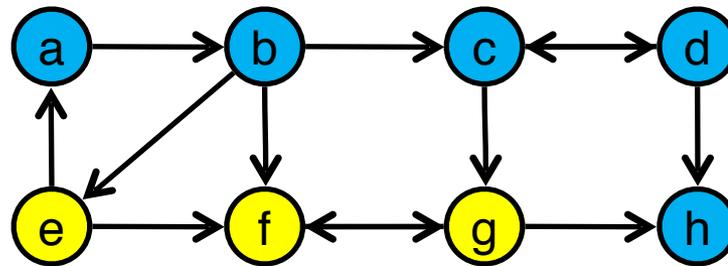


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SCC

- Algorithm

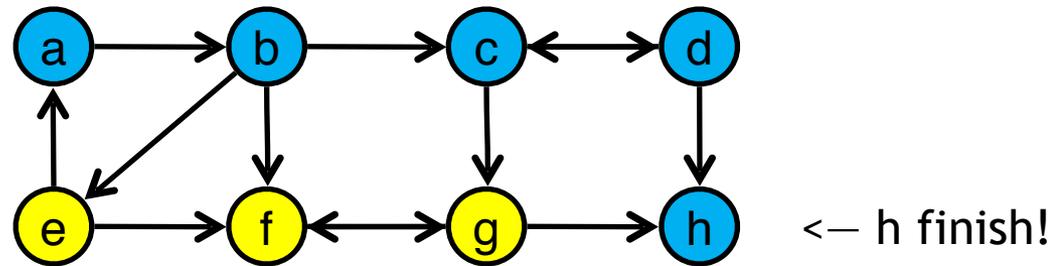


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SCC

- Algorithm

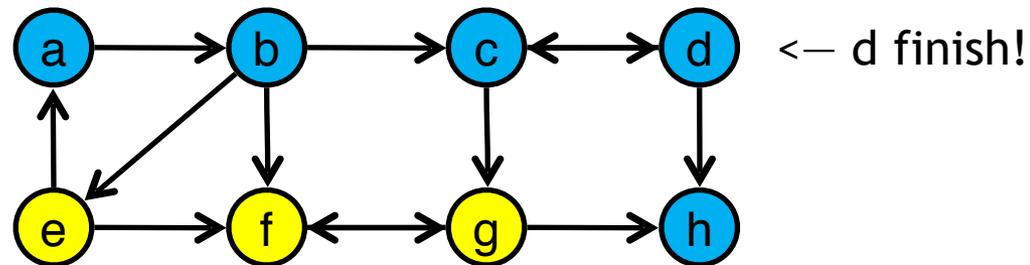


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SCC

- Algorithm

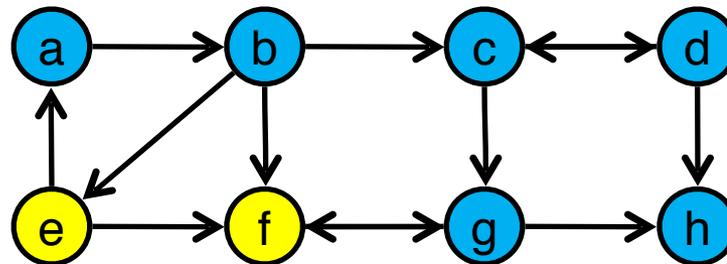


h	d								
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SCC

- Algorithm

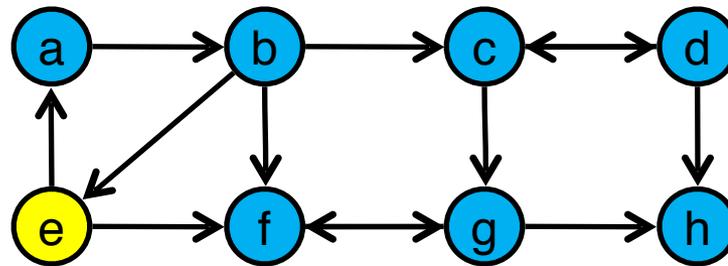


h	d								
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SCC

- Algorithm

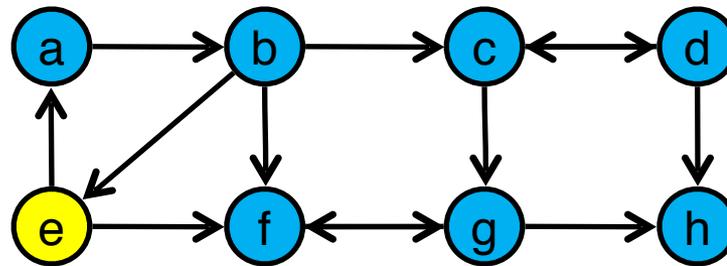


h	d								
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SCC

- Algorithm



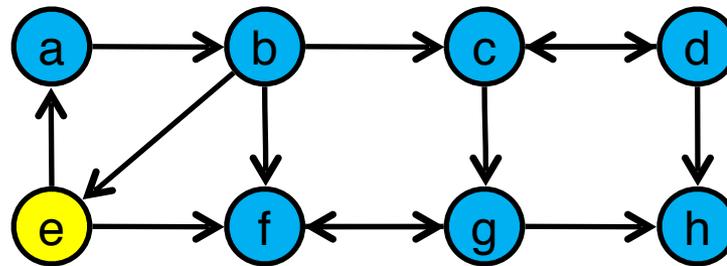
^ f finish!

h	d	f							
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SCC

- Algorithm



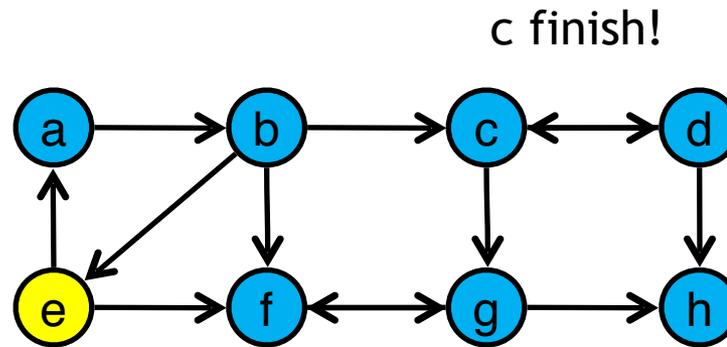
^ g finish!

h	d	f	g						
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SCC

- Algorithm

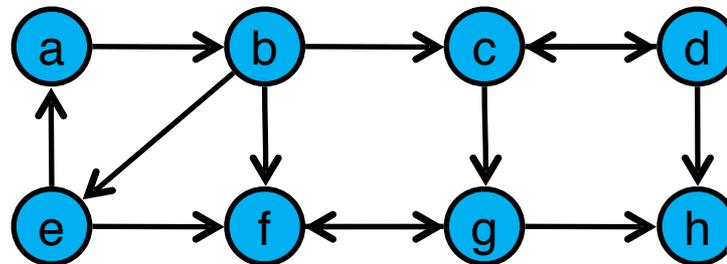


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SCC

- Algorithm

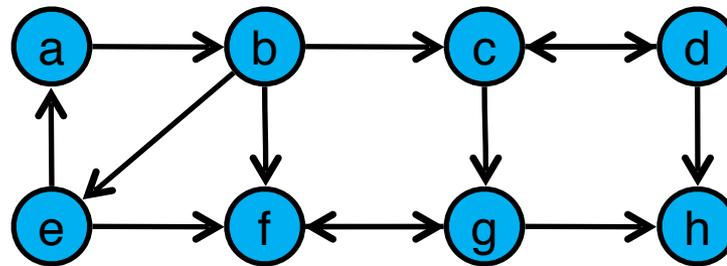


h	d	f	g	c					
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SCC

- Algorithm



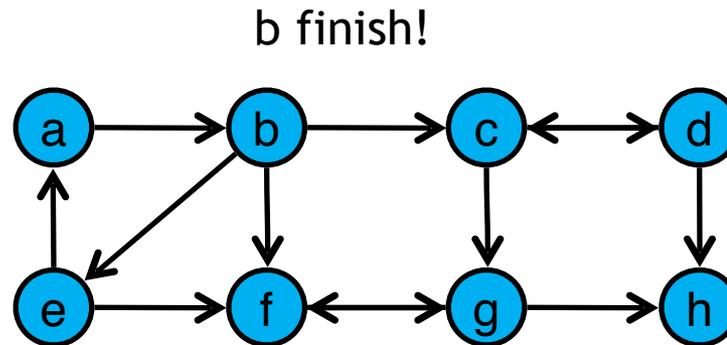
^ e finish!

h	d	f	g	c	e				
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SCC

- Algorithm



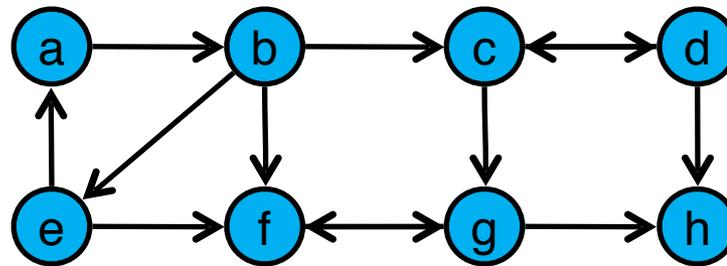
h	d	f	g	c	e	b			
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SCC

- Algorithm

a finish!

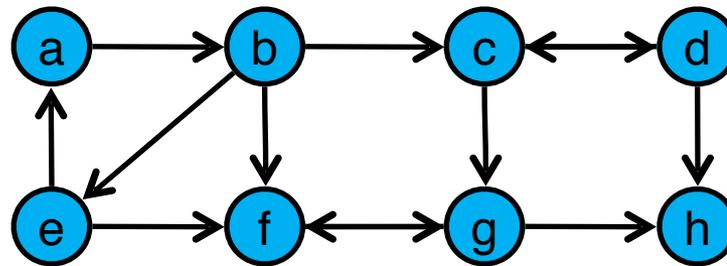


h	d	f	g	c	e	b	a		
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SCC

- Algorithm
 - Reverse the graph

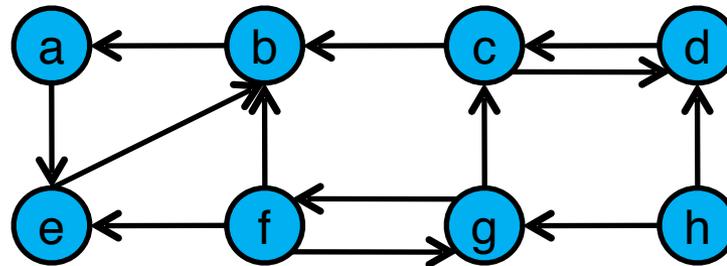


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SCC

- Algorithm
 - Reverse the graph

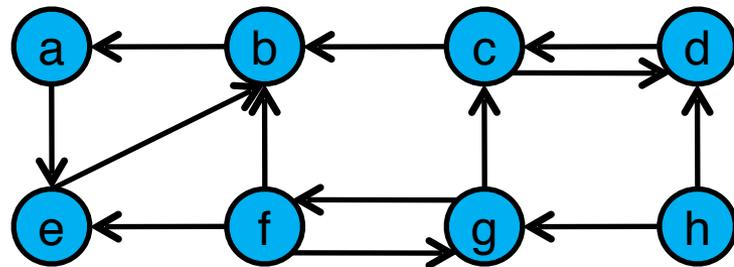


h	d	f	g	c	e	b	a		
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SCC

- Algorithm
 - Reverse the graph
 - Re-search by the ending time

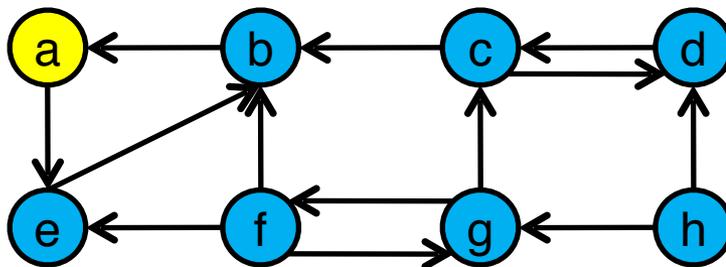


h	d	f	g	c	e	b	a		
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SCC

- Algorithm
 - Reverse the graph
 - Re-search by the ending time

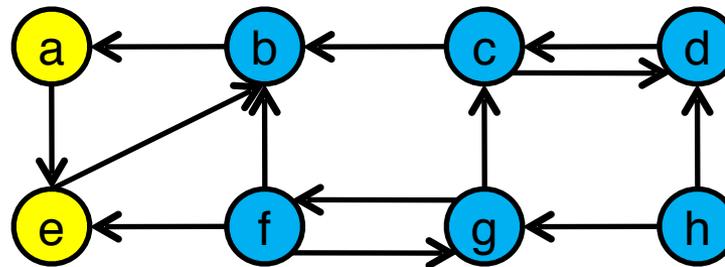


h	d	f	g	c	e	b	a		
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SCC

- Algorithm
 - Reverse the graph
 - Re-search by the ending time

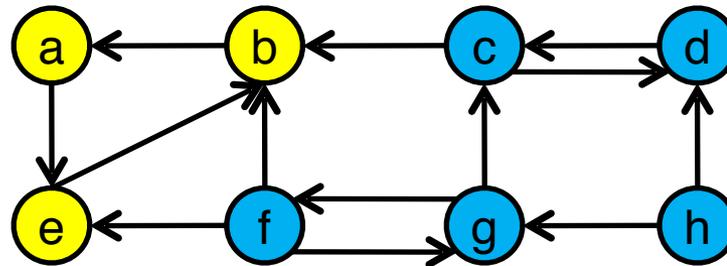


h	d	f	g	c	e	b	a		
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SCC

- Algorithm
 - Reverse the graph
 - Re-search by the ending time

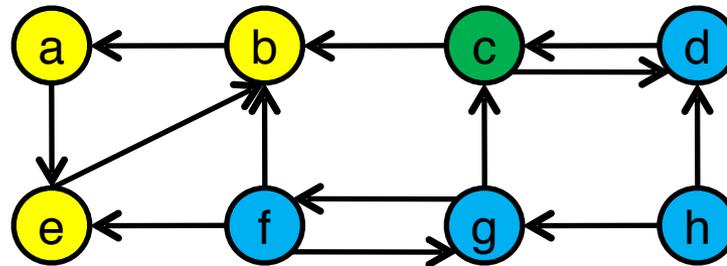


h	d	f	g	c	e	b	a		
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SCC

- Algorithm
 - Reverse the graph
 - Re-search by the ending time

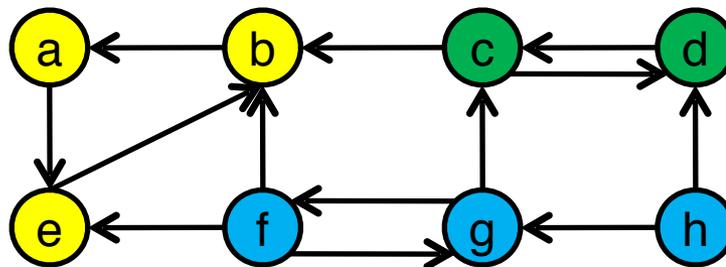


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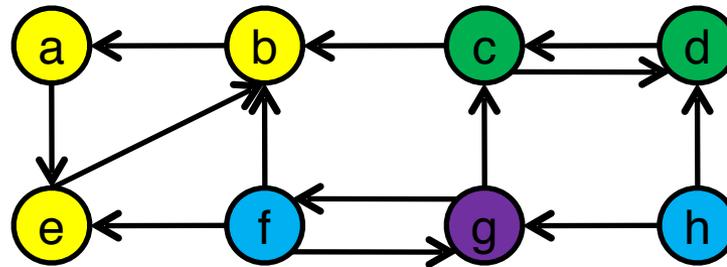
SCC

- Algorithm
 - Reverse the graph
 - Re-search by the ending time



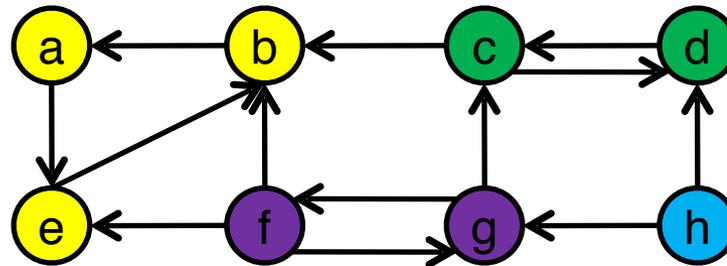
SCC

- Algorithm
 - Reverse the graph
 - Re-search by the ending time



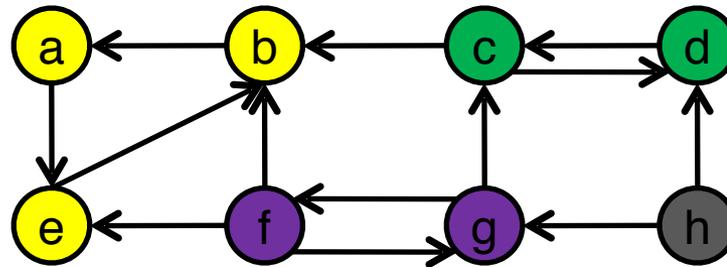
SCC

- Algorithm
 - Reverse the graph
 - Re-search by the ending time



SCC

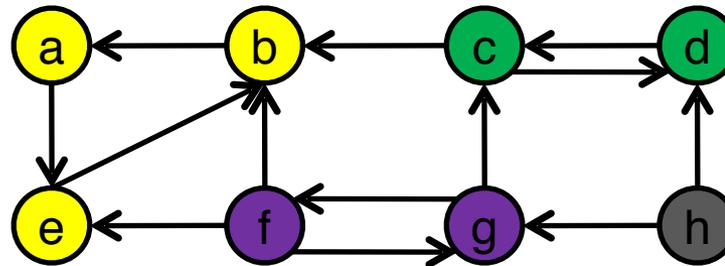
- Algorithm
 - Reverse the graph
 - Re-search by the ending time



SCC

- Algorithm
 - Reverse the graph
 - Re-search by the ending time

4 components



Practice

ICPC - 4262



Learn more

- Tarjan's algorithm
 - only one DFS



Homework

- UVA (total 14 problems)
 - 247, 315, 459, 610, 796, 10199, 10731, 10765, 11324, 11504, 11709, 11770, 11838, 12783
- POJ (total 5 problems)
 - 1236, 1523, 2117, 2186, 2553
- ICPC (total 3 problems)
 - 4262, 4839, 5135

基本門檻 5 題，第二次修課同學請從橘色的題號選擇

