

# NCKU Programming Contest Training Course

## Course 9

### 2015/03/25

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# Outline

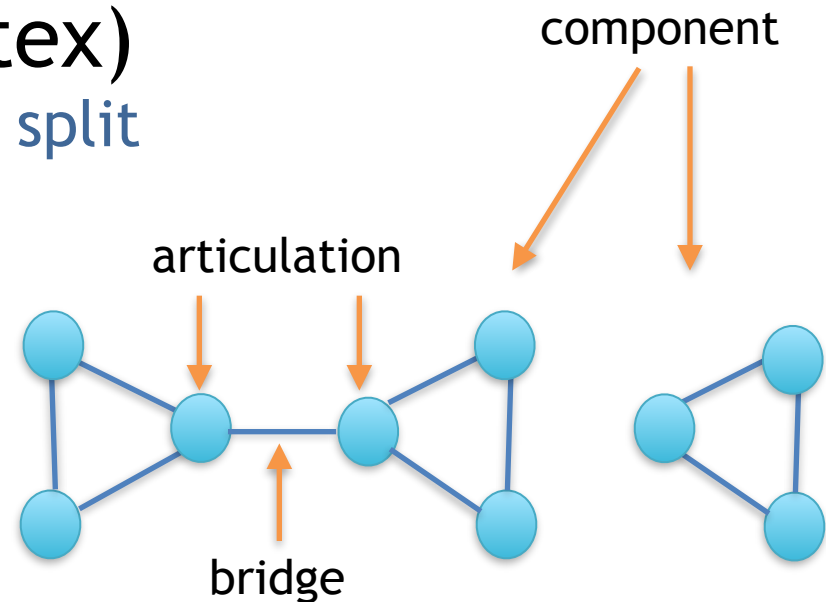
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- Articulation/Bridge
- Strongly Connected Component(SCC)



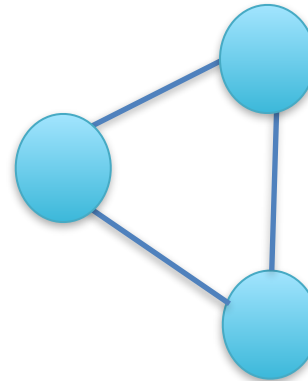
# Graph

- **connected graph/component**  
connected graph/component iff all pairwise vertices exist **at least one path** & **no more vertices can be added**
- **articulation(cut-vertex)**  
remove articulation vertex split one component to two
- **bridge(cut-edge)**  
same as articulation



# Articulation/Bridge

- Find Articulation in Graph
  - Graph become **non-connected** if remove a Articulation.  
 $V$  times DFS =  $O(V*(V+E)) \rightarrow$  **too slow!**
  - Vertex is not Articulation if can find **alternative path**  
 $\rightarrow$  find cycle!
  - Use DFS  $\rightarrow O(V+E)$



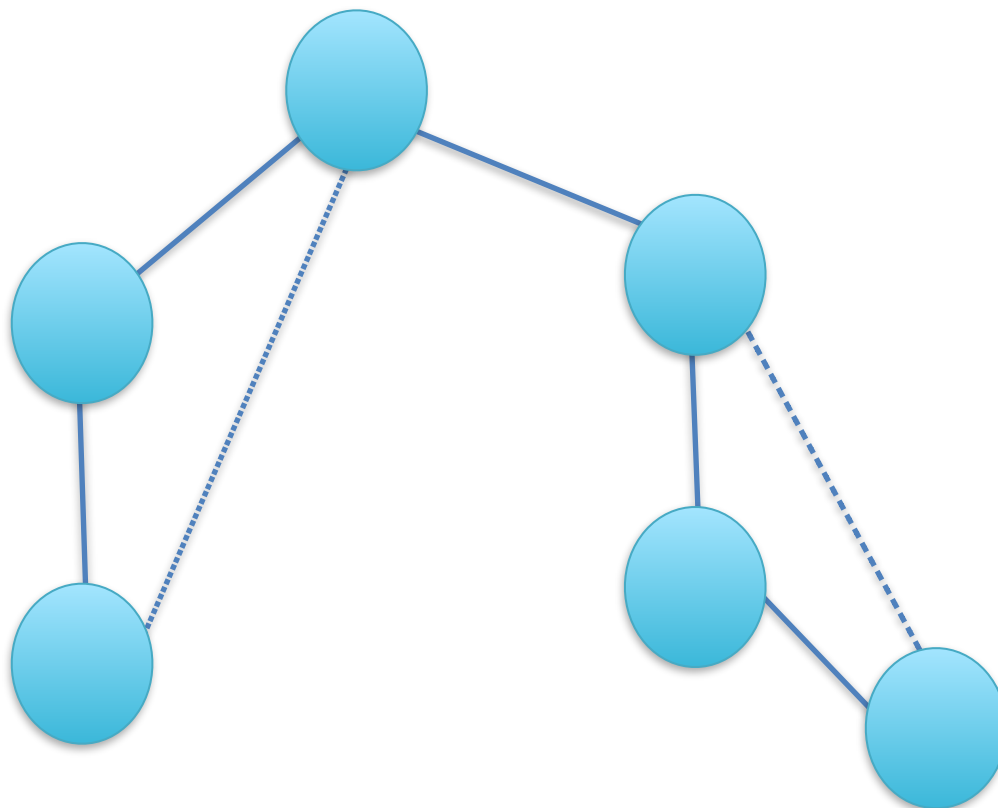
# Articulation/Bridge

- Concept
  - if vertex  $u$ 's children can't back to  $u$ 's ancestors  
→  $u$  is Articulation
  - if vertex  $u$  is root and has at least 2 child  
→  $u$  is Articulation
- Bridge?
  - two Articulation  $u, v$  have an edge →  $(u, v)$  is Bridge!

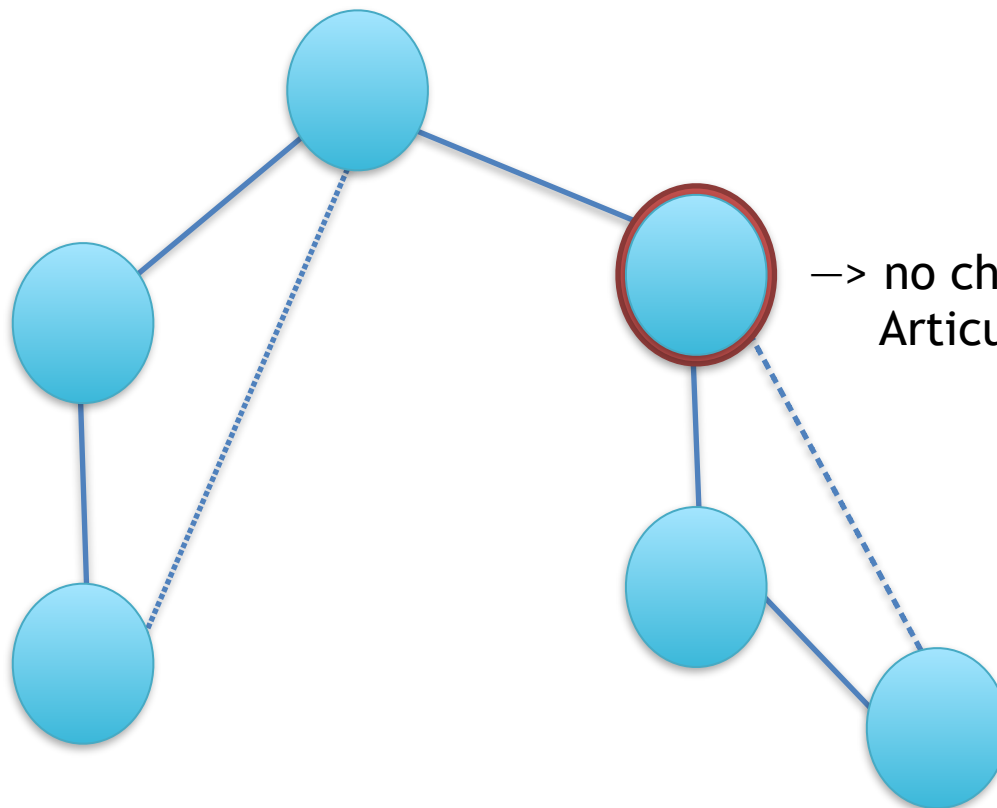


# Articulation/Bridge

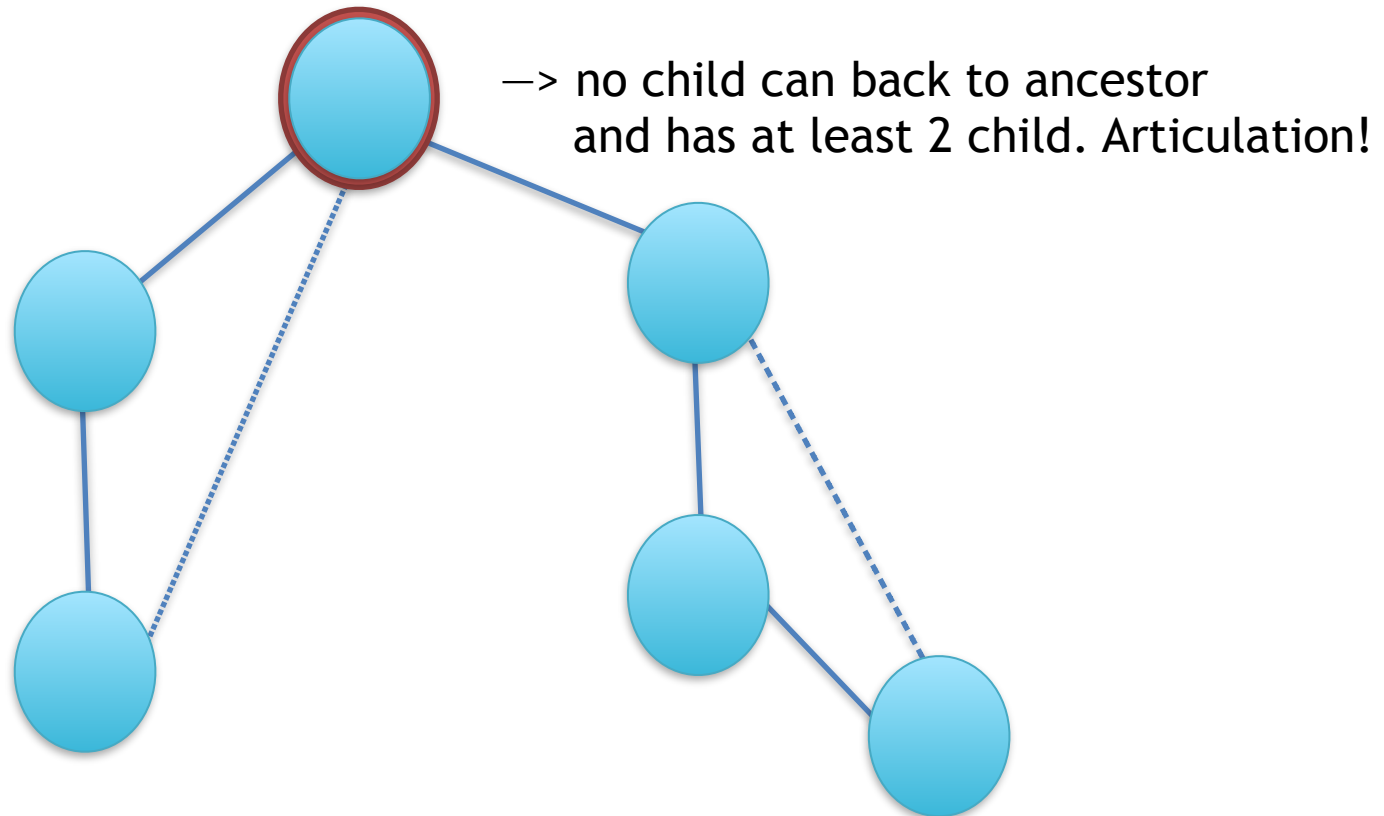
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# Articulation/Bridge

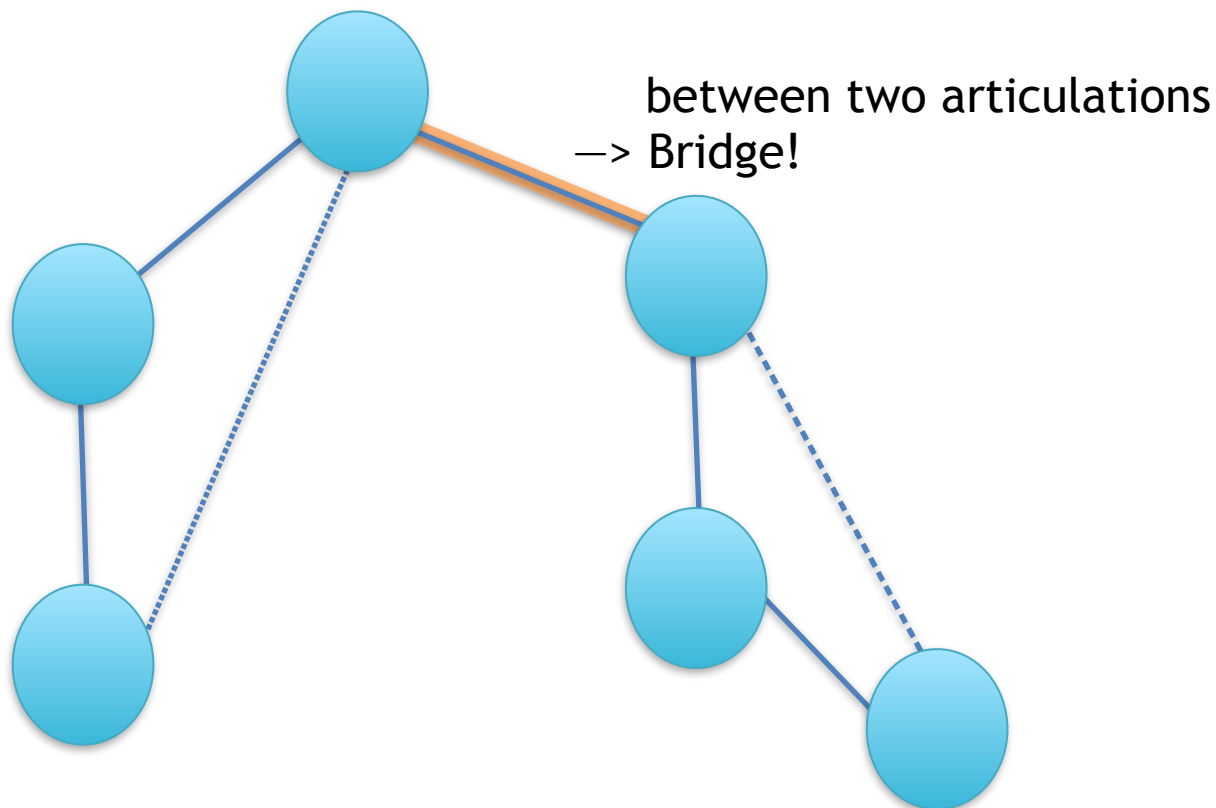


# Articulation/Bridge



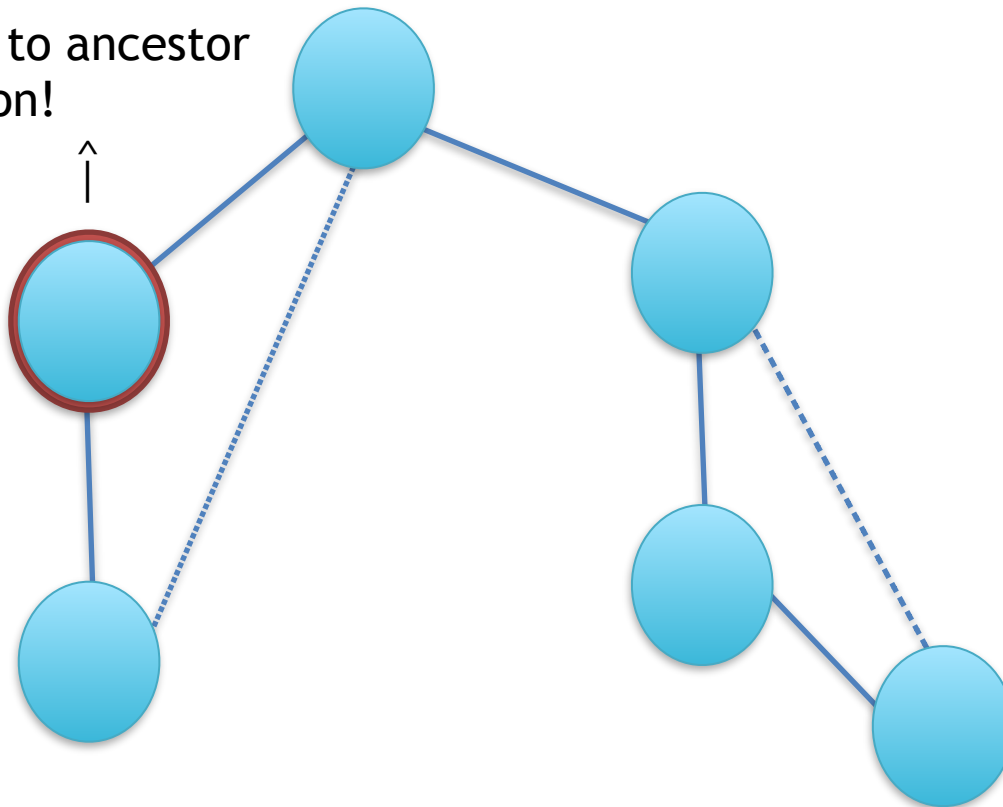


# Articulation/Bridge



# Articulation/Bridge

child can back to ancestor  
NOT Articulation!



# Articulation / Bridge

- $dfn[u]$  = DFS traversal order
  - first visit time each vertex  $u$  in DFS
- $low[u] = \min(dfn[u], \text{lowest } low[v])$ 
  - if edge  $(u,v)$  exist and  $v$  is not  $u$ 's parent



# Articulation / Bridge

- Articulation
  - if vertex  $u$ 's children can't back to  $u$ 's ancestors  
→  $dfn[u] \leq low[v]$ ,  $v$  is  $u$ 's child
  - if vertex  $u$  is root and has at least 2 child  
→ count child  $\geq 2$
- Bridge?
  - two Articulation  $u, v$  →  $dfn[u] < low[v]$ ,  $v$  is  $u$ 's child



# Practice

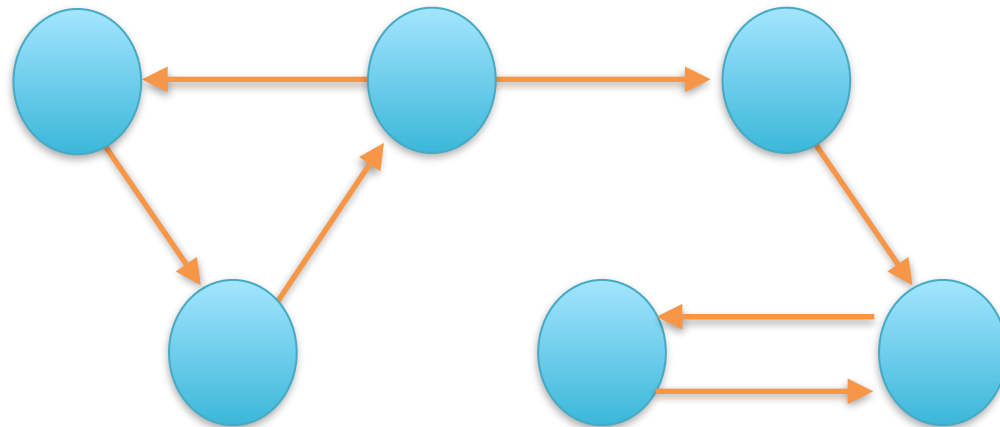
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## UVA - 315

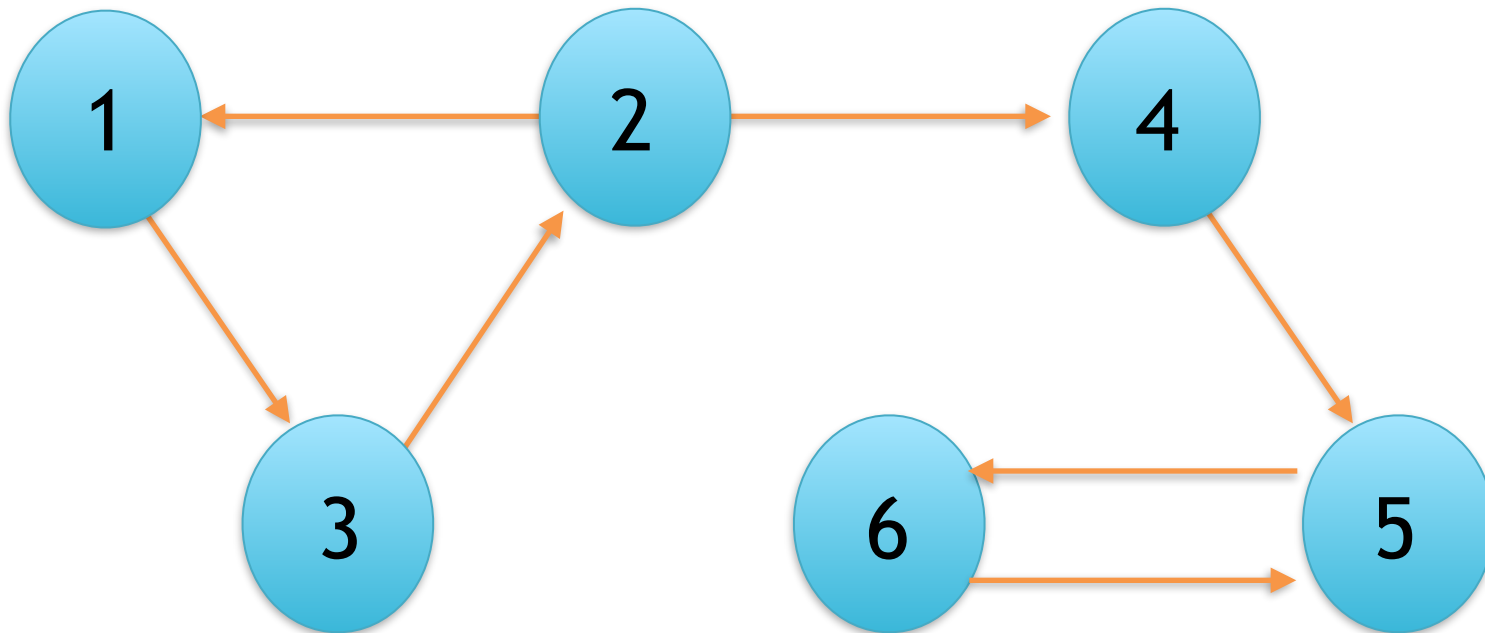


# SCC

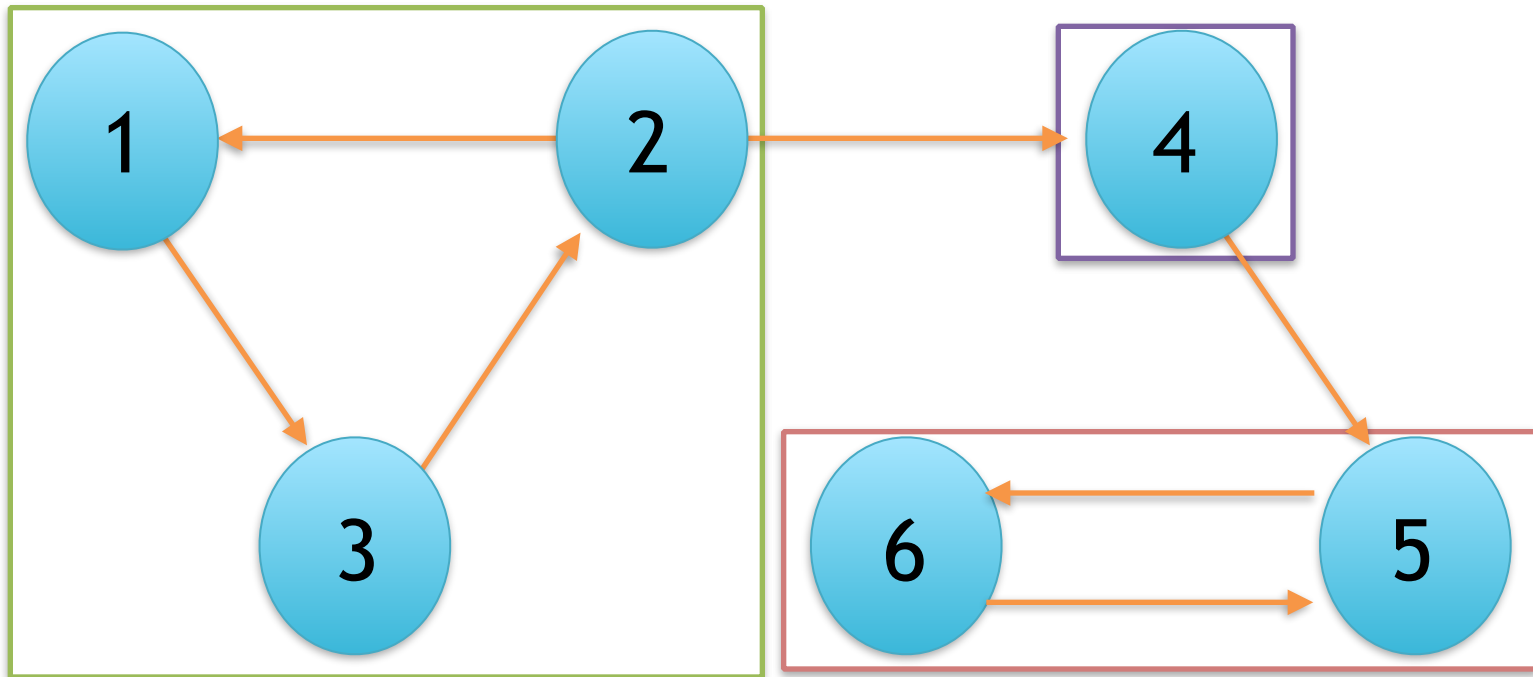
- connected component in **directed** graph
  - same definition in undirected graph



# SCC



# SCC

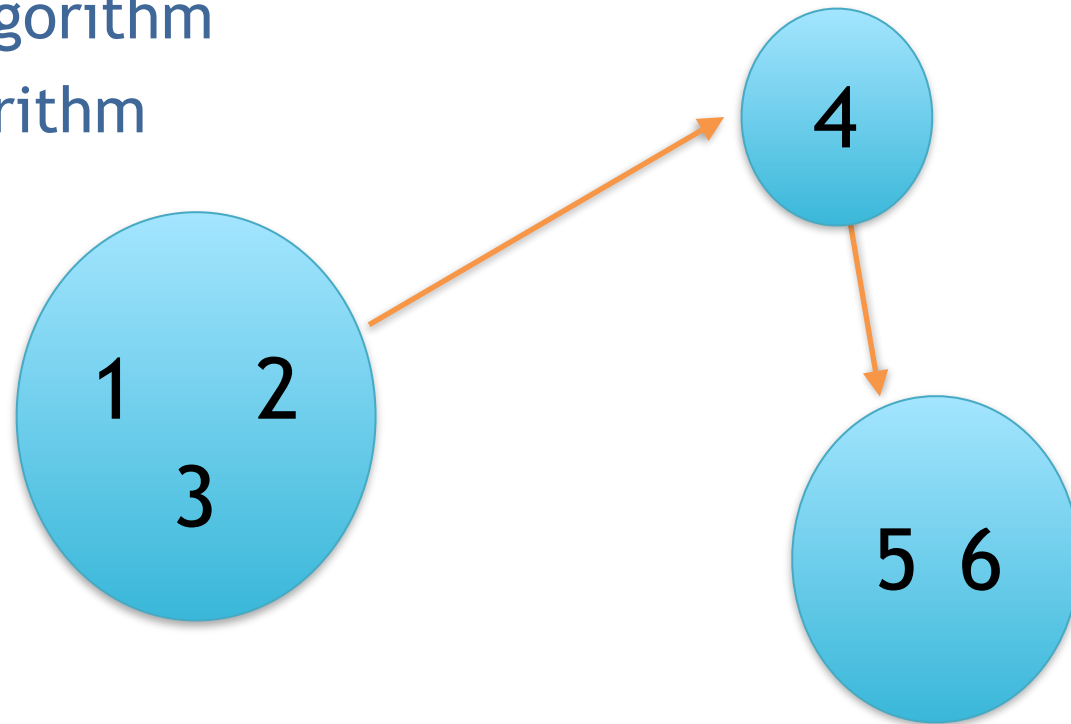




# SCC

find all SCCs, **contract all cycles** → DAG (directed acyclic graph)

- Kosaraju's Algorithm
- Tarjan's Algorithm



# SCC

- Kosaraju's algorithm

## STRONGLY-CONNECTED-COMPONENTS( $G$ )

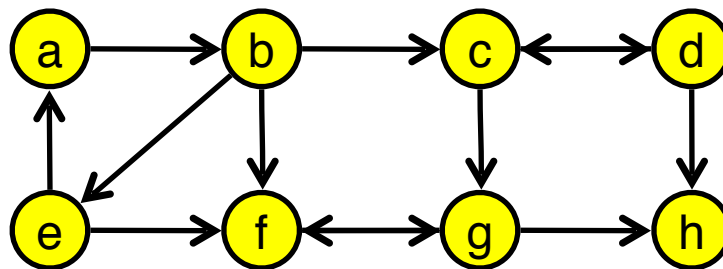
1. Call DFS( $G$ ) to compute finishing time for each vertex.
2. Compute transpose of  $G$  i.e.,  $G^T$ .
3. Call DFS( $G^T$ ) but this time consider the vertices in order of decreasing finish time.
4. Out the vertices of each tree in DFS-forest.

twice DFS  $\rightarrow$  total complexity:  $O(V+E)$



# SCC

- Algorithm

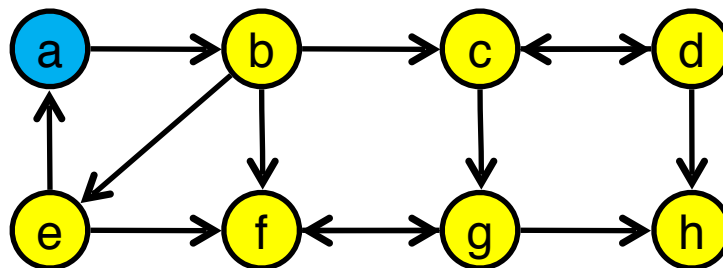


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# SCC

- Algorithm

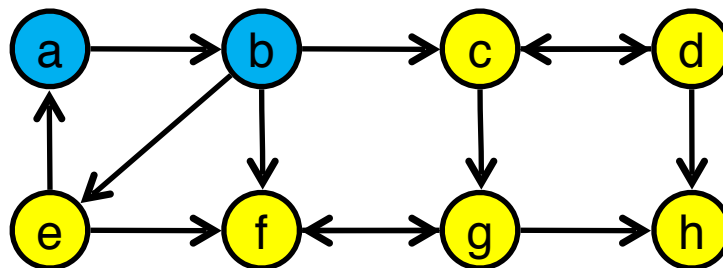


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# SCC

- Algorithm

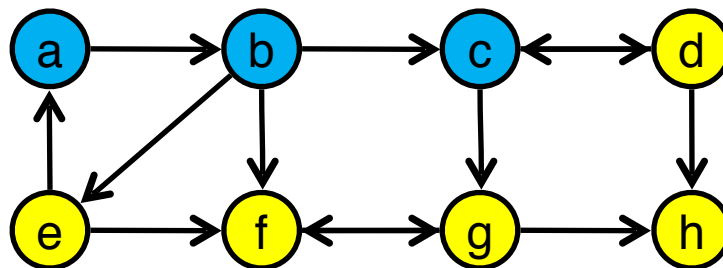


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# SCC

- Algorithm

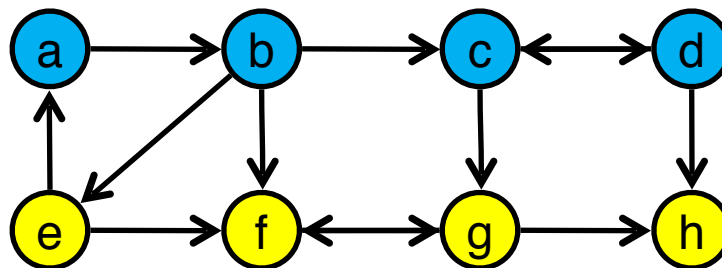


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# SCC

- Algorithm

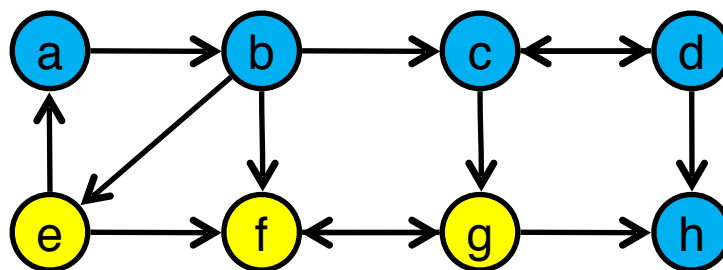


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# SCC

- Algorithm

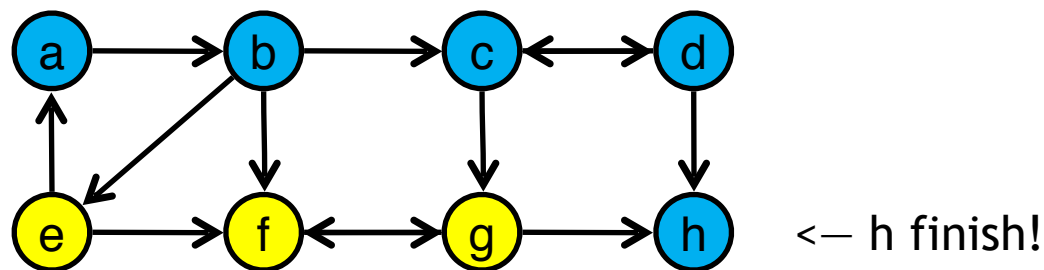


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# SCC

- Algorithm

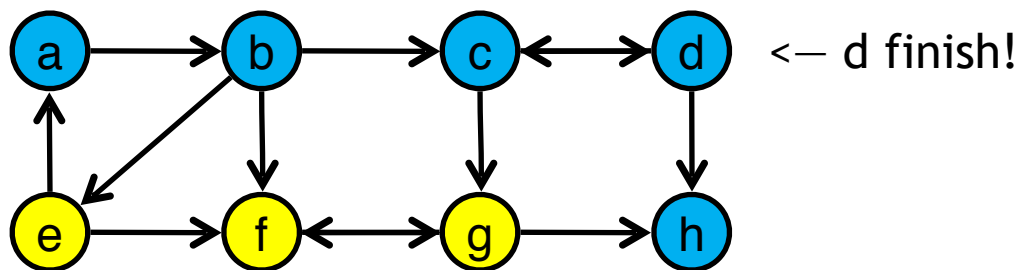


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# SCC

- Algorithm

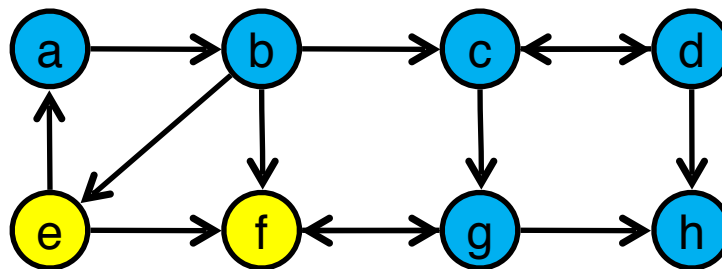


h	d								
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# SCC

- Algorithm

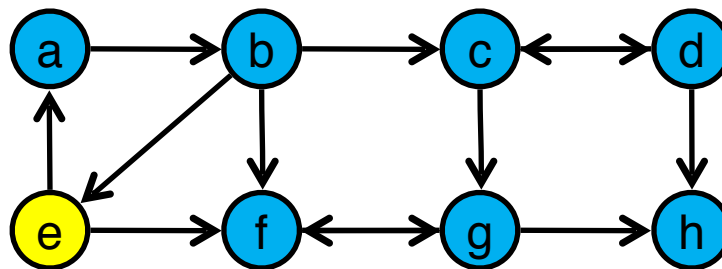


h	d								
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# SCC

- Algorithm

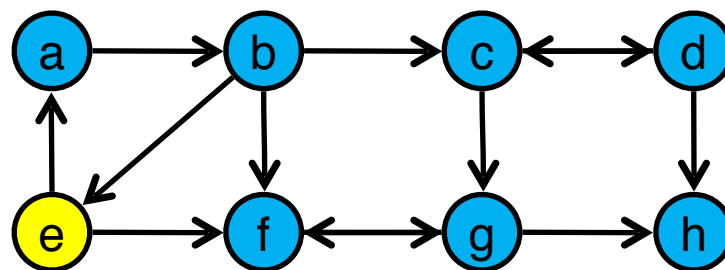


h	d								
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# SCC

- Algorithm



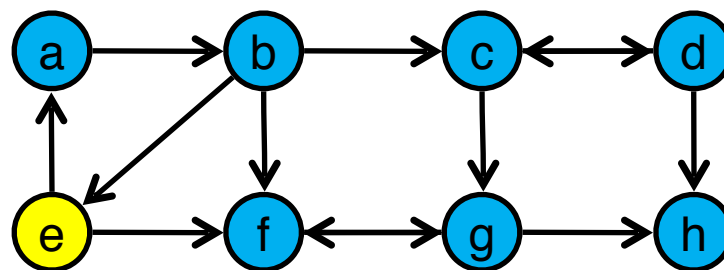
^ f finish!

h	d	f							
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# SCC

- Algorithm



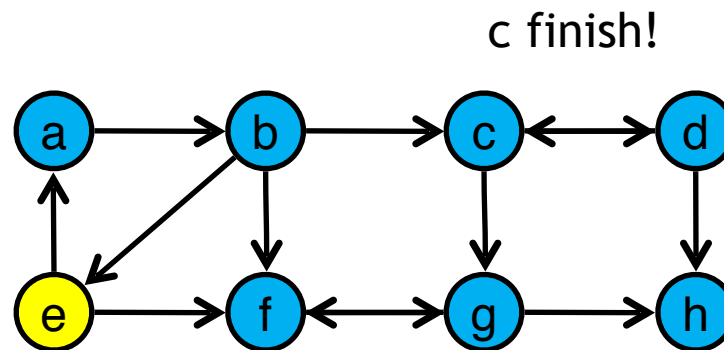
^ g finish!

h	d	f	g						
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# SCC

- Algorithm

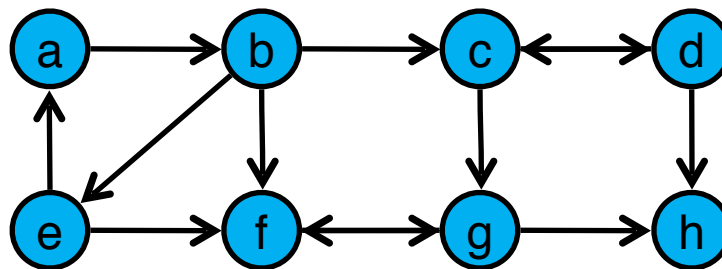


h	d	f	g	c					
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# SCC

- Algorithm



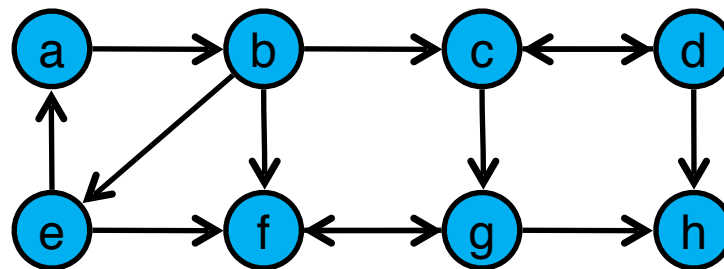
h	d	f	g	c					
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# SCC

- Algorithm



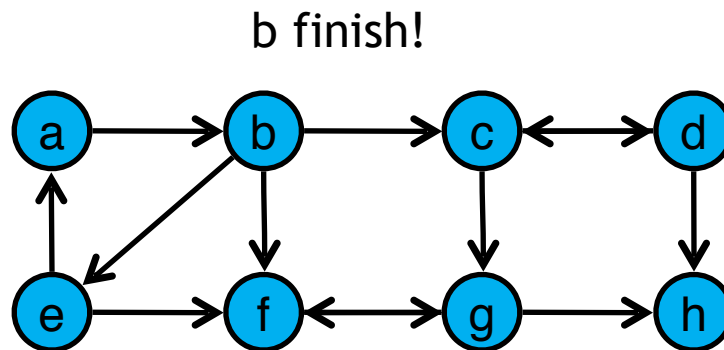
^ e finish!

h	d	f	g	c	e				
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# SCC

- Algorithm



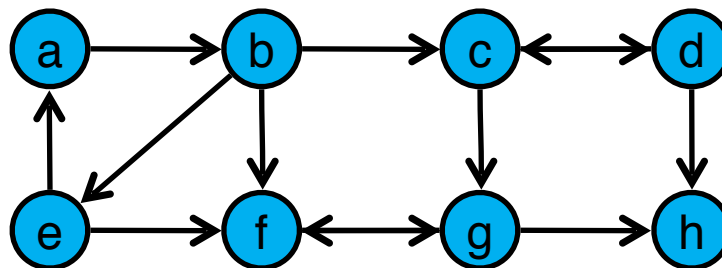
h	d	f	g	c	e	b			
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# SCC

- Algorithm

a finish!

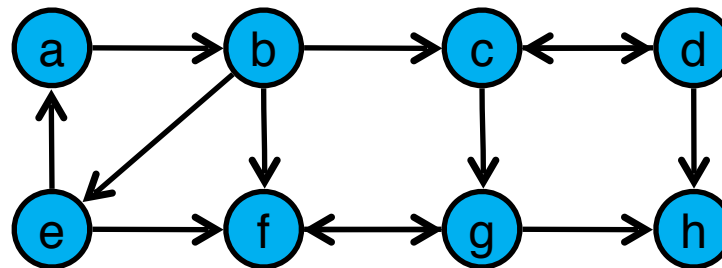


h	d	f	g	c	e	b	a		
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# SCC

- Algorithm
  - Reverse the graph

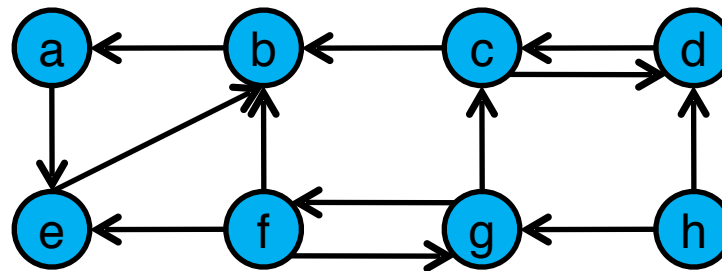


h	d	f	g	c	e	b	a		
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# SCC

- Algorithm
  - Reverse the graph

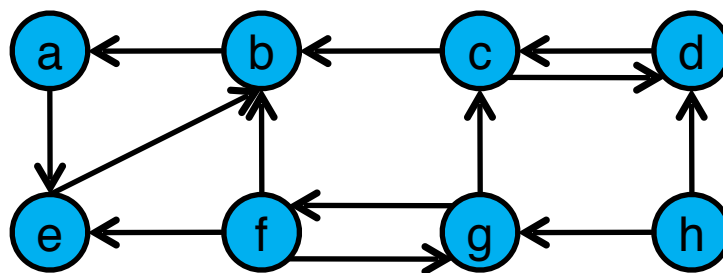


h	d	f	g	c	e	b	a		
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# SCC

- Algorithm
  - Reverse the graph
  - Re-search by the ending time

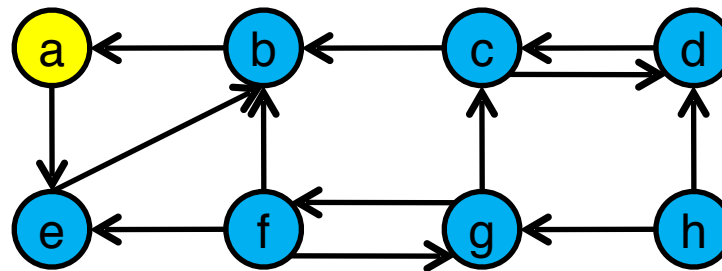


h	d	f	g	c	e	b	a		
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# SCC

- Algorithm
  - Reverse the graph
  - Re-search by the ending time

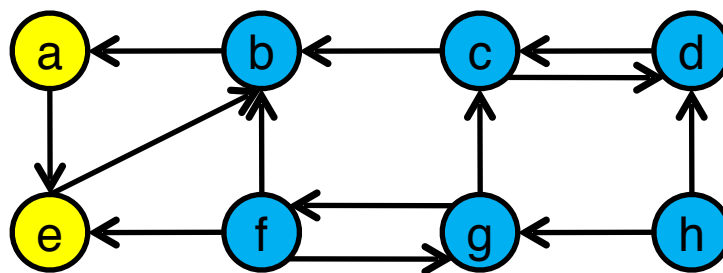


h	d	f	g	c	e	b	a		
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# SCC

- Algorithm
  - Reverse the graph
  - Re-search by the ending time



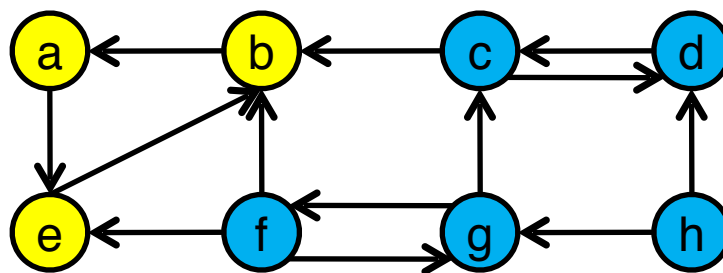
h	d	f	g	c	e	b	a		
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# SCC

- Algorithm
  - Reverse the graph
  - Re-search by the ending time

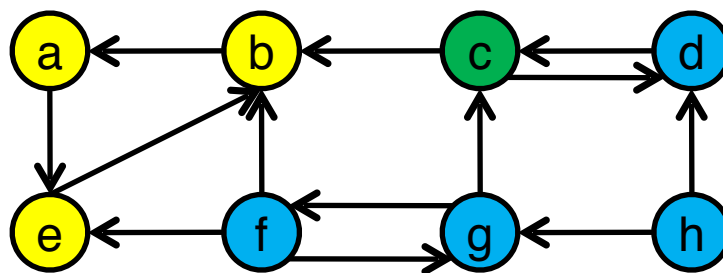


h	d	f	g	c	e	b	a		
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# SCC

- Algorithm
  - Reverse the graph
  - Re-search by the ending time

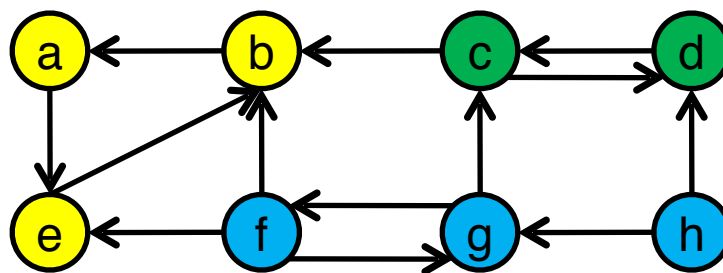


h	d	f	g	c	e	b	a		
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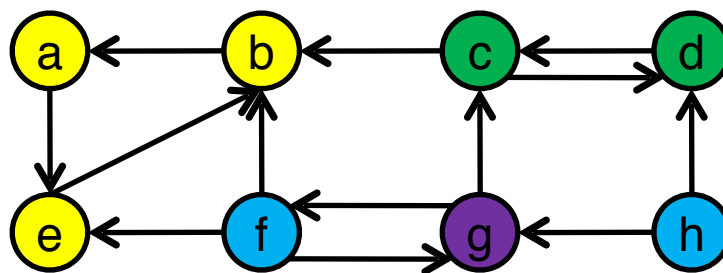
# SCC

- Algorithm
  - Reverse the graph
  - Re-search by the ending time



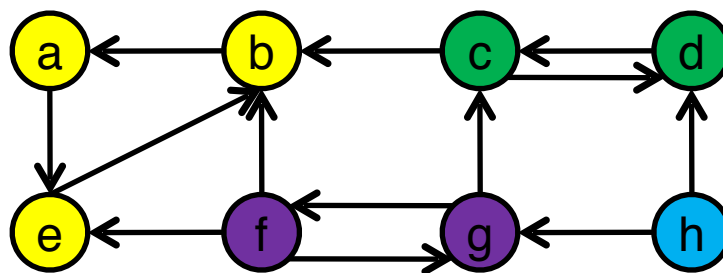
# SCC

- Algorithm
  - Reverse the graph
  - Re-search by the ending time



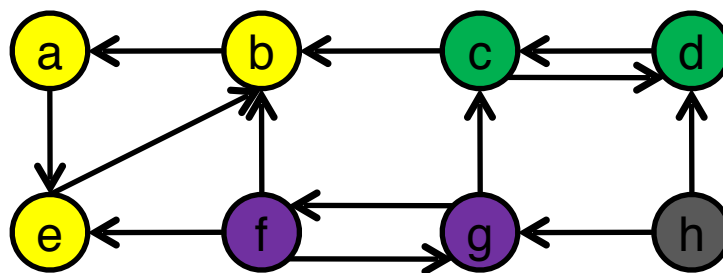
# SCC

- Algorithm
  - Reverse the graph
  - Re-search by the ending time



# SCC

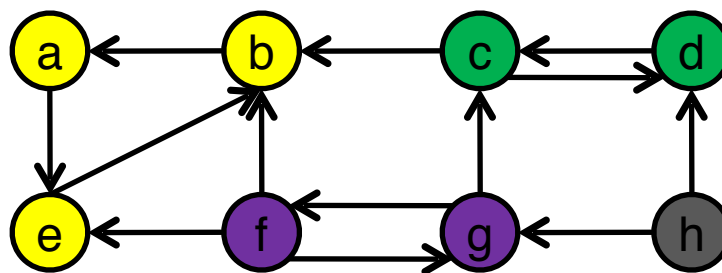
- Algorithm
  - Reverse the graph
  - Re-search by the ending time



# SCC

- Algorithm
  - Reverse the graph
  - Re-search by the ending time

4 components



# Practice

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## ICPC - 4262





# Learn more

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- Tarjan's algorithm
  - only one DFS



# Homework

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- UVA (total 14 problems)
  - 247, 315, 459, 610, 796, 10199, 10731, 10765, 11324, 11504, 11709, 11770, 11838, 12783
- POJ (total 5 problems)
  - 1236, 1523, 2117, 2186, 2553
- ICPC (total 3 problems)
  - 4262, 4839, 5135

基本門檻 5 題，第二次修課同學請從橘色的題號選擇

